

FIG. 1

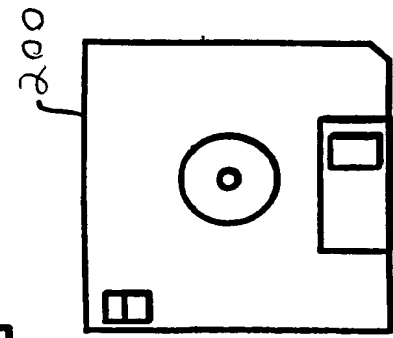


FIG. 2

```
<PurchaseOrder>
  <Buyer>
    <Name>John Smith</Name>
  </Buyer>
  <LineItems>
    <Item>
      <Name>ThinkPad T20</Name>
      <Description>
        This laptop is the <Bold>best value</Bold> in its class
      </Description>
      <:Price>$1999</Price>
    </Item>
    <Item>
      <Name>NetVista Desktop</Name>
      <:Price>$999</Price>
    </Item>
  </LineItems>
</PurchaseOrder>
```

Figure 3

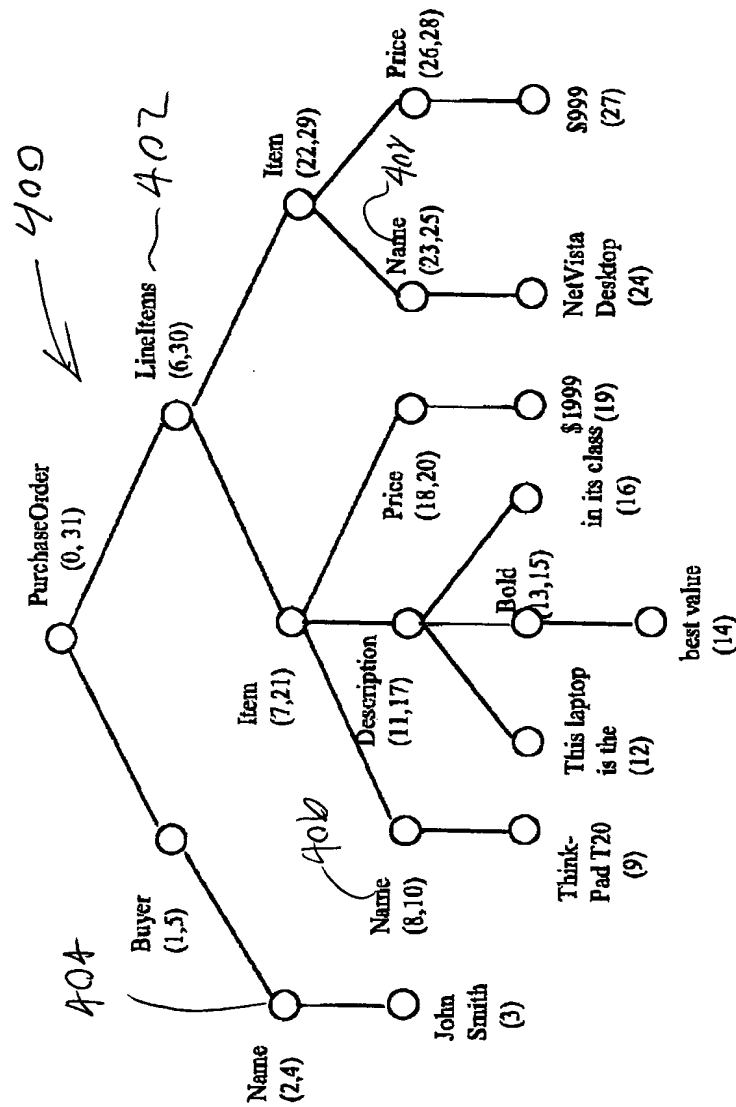
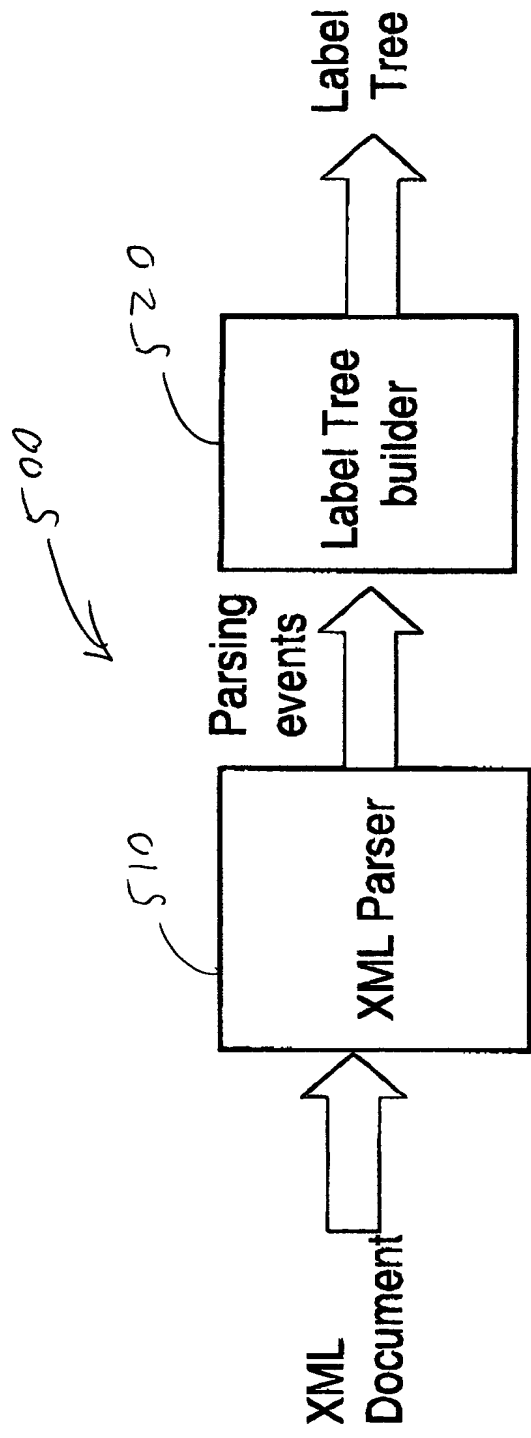


FIG. 4



Building the label tree in bulk-loading mode

FIG. 5

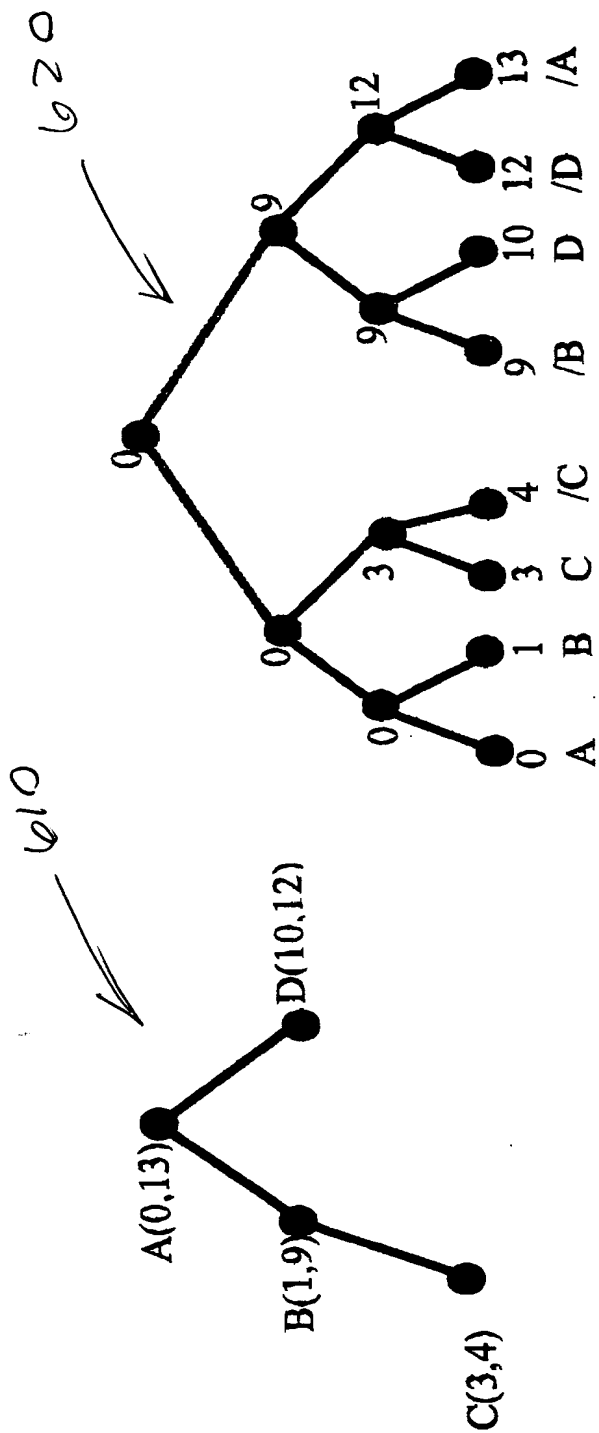
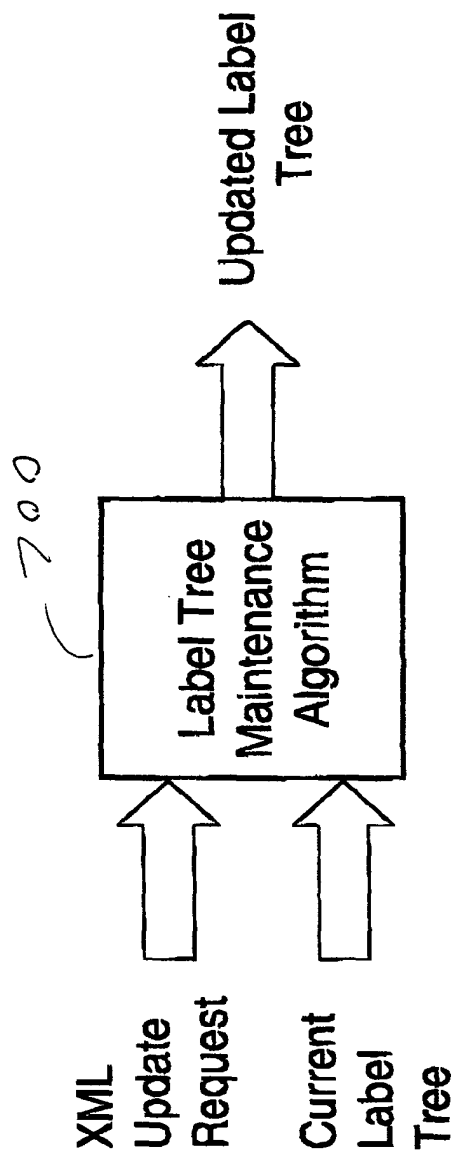


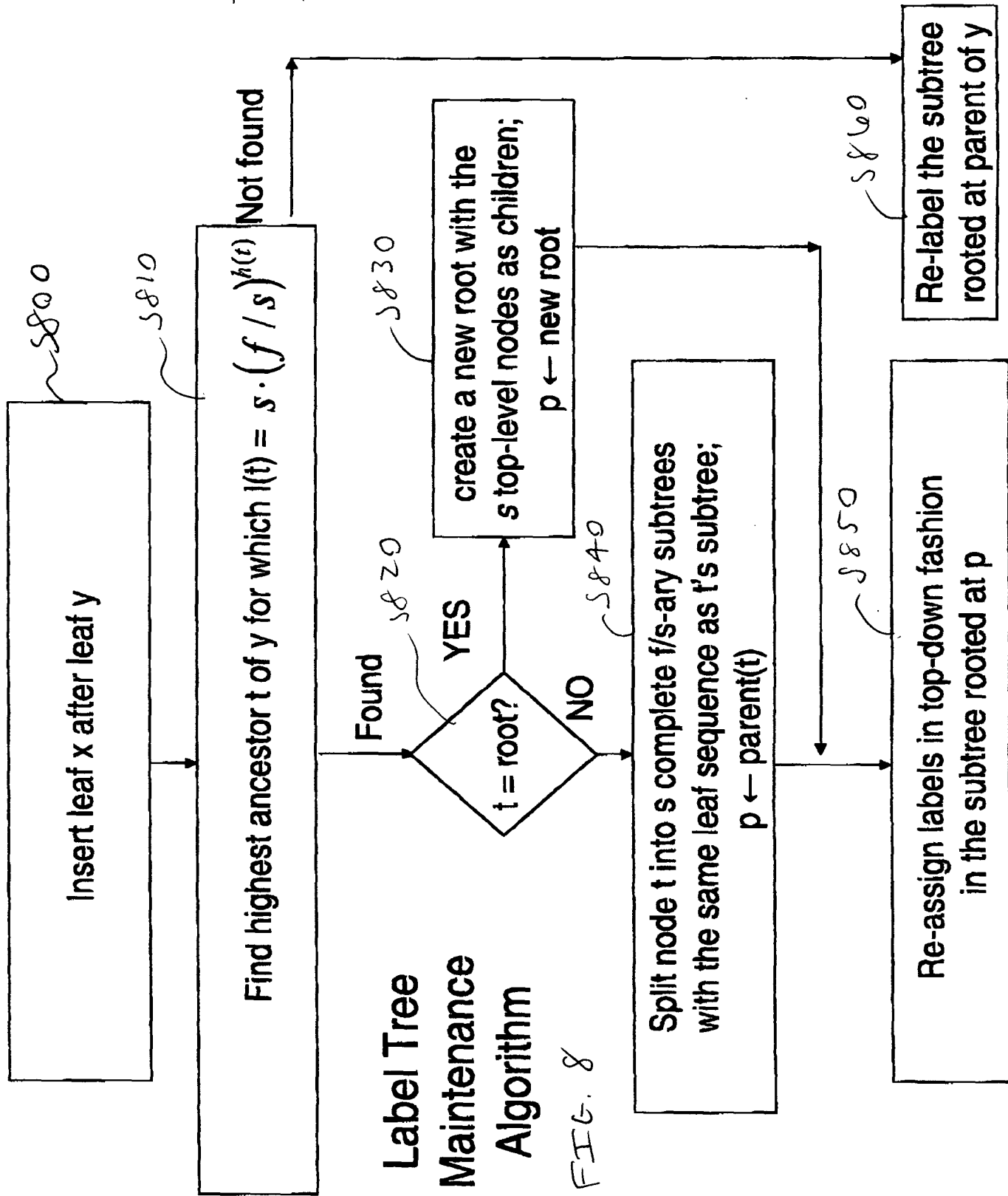
FIG. 6B

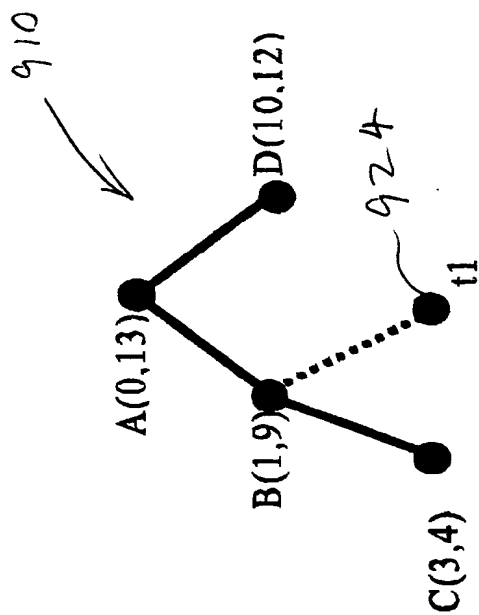
FIG. 6A



Maintaining the label tree in response to updates

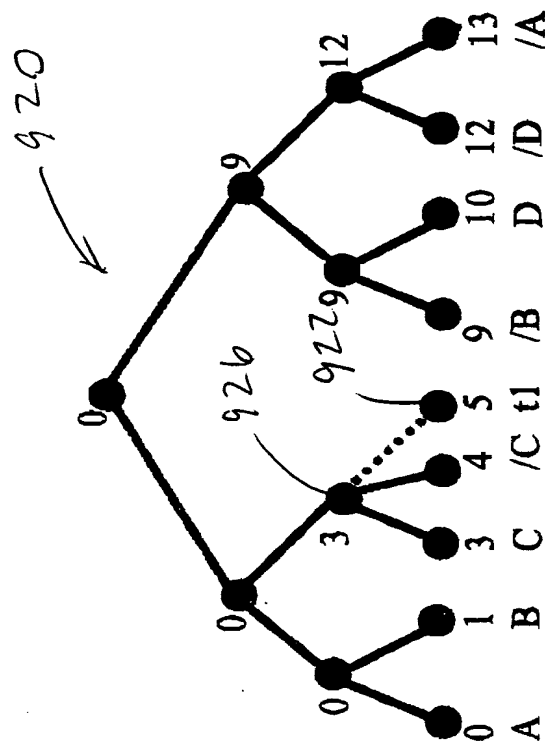
FIG. 7





XML tree

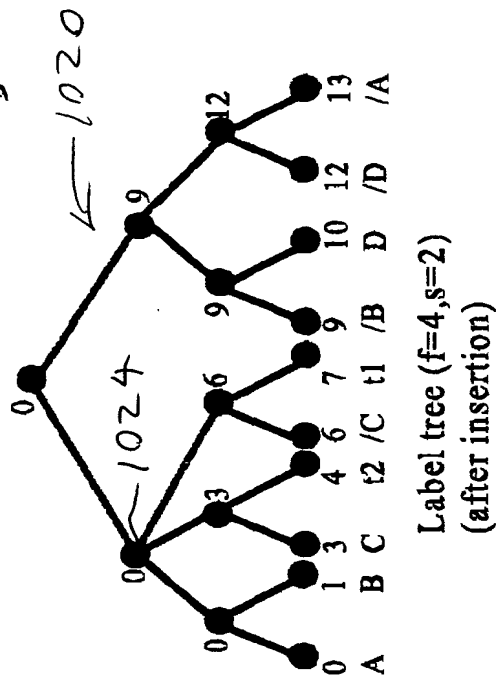
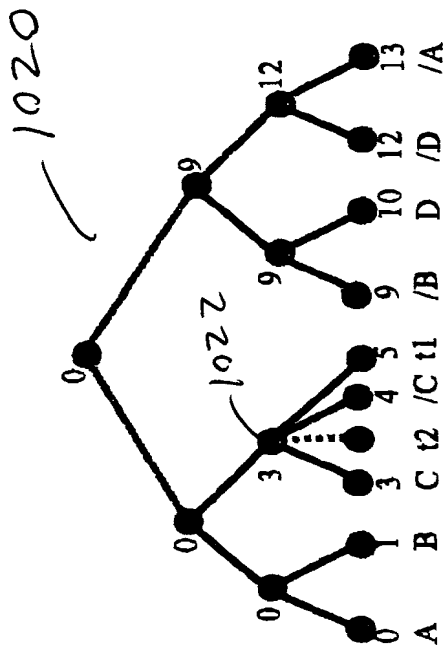
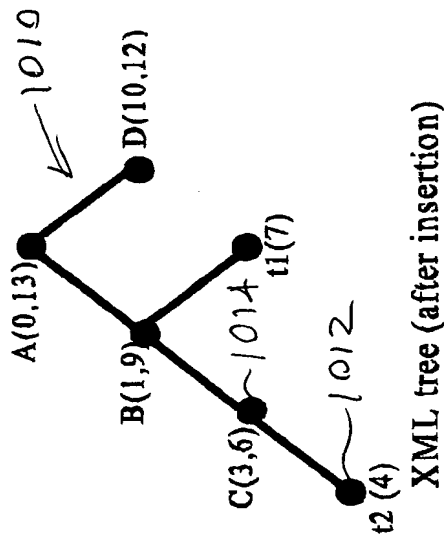
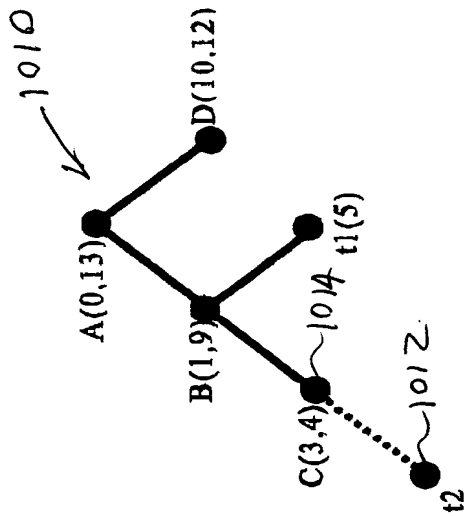
FIG 9A



Label tree (f=4,s=2)

FIG 9B





Amortized cost upper bound  
( $f=4, s=2$ )

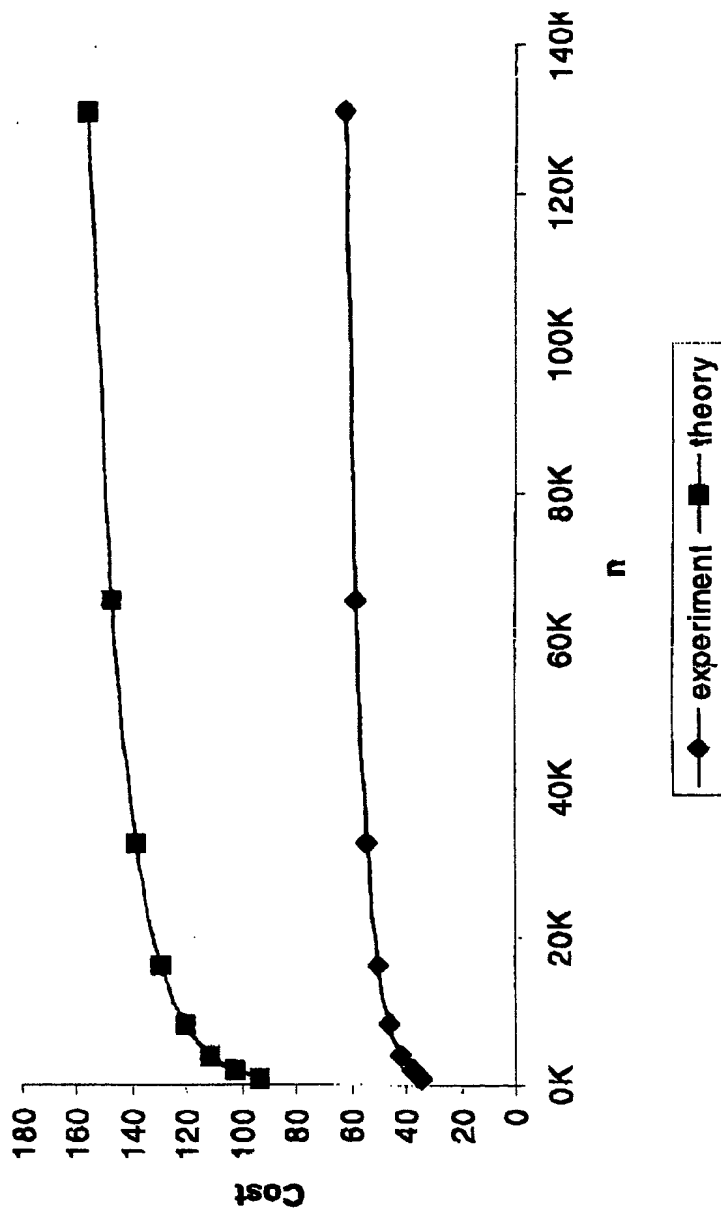


Fig. 11

Amortized cost by varying  $f$   
 $(n=2^{15}, s=4)$

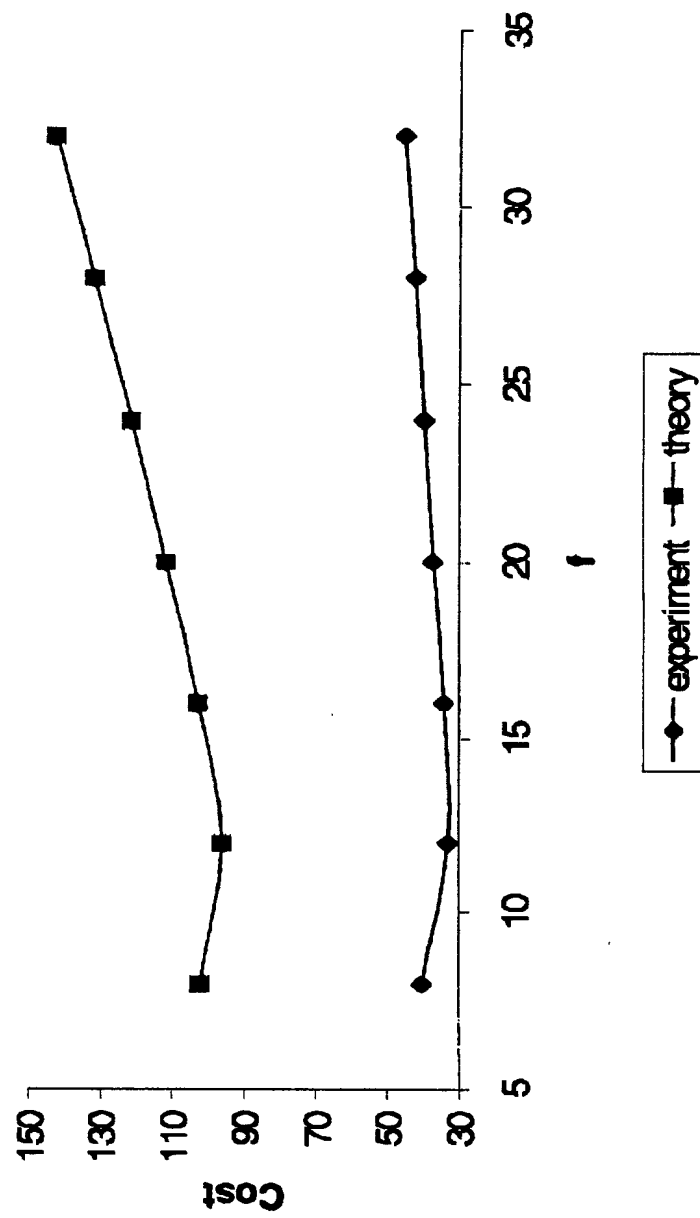


Fig. 12

Amortized cost by varying  $s$   
 $(n=2^{15}, f=36)$

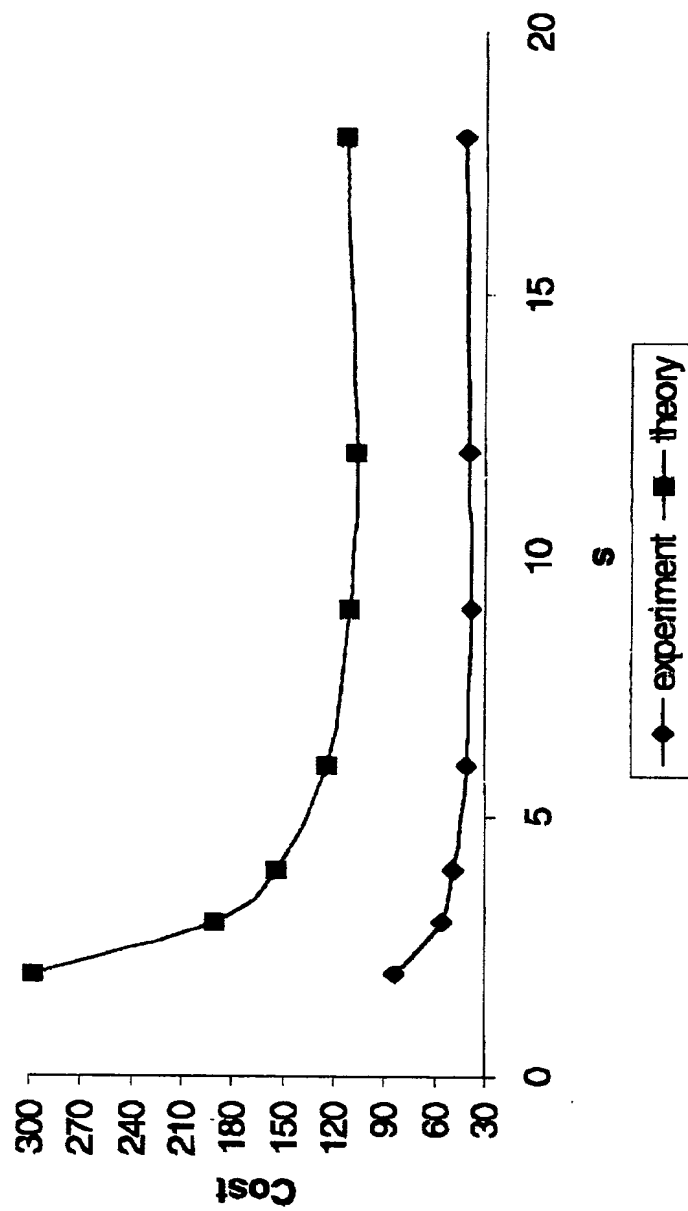


FIG. 13

**Optimal cost constrained on bits used  
( $n=2^{15}$ )**

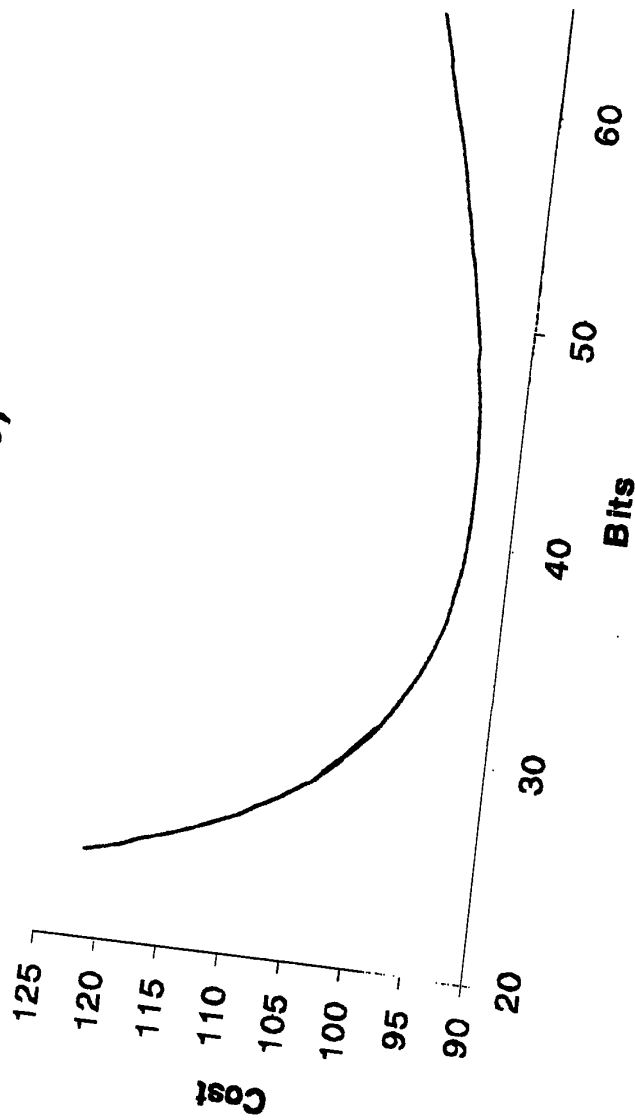


FIG. 14